

50th Anniversary of the Birth of Prolog:

Some reflections on Prolog's Evolution, Status, and Future

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Part of the contents of this talk appear in the recent TPLP paper “50 years of Prolog and Beyond,” by

*Philipp Körner, Michael Leuschel, João Barbosa, Vítor Santos Costa, Verónica Dahl,
Manuel V. Hermenegildo, Jose F. Morales, Jan Wielemaker, Daniel Diaz,
Salvador Abreu, and Giovanni Ciatto*

written for Prolog's 50th anniversary and TPLP's 20th anniversary.

The Year of Prolog

- Summer of 1972:
Alain Colmerauer and team in Marseille develop the first version of Prolog.
- This event + earlier and later collaborations w/Bob Kowalski and colleagues in Edinburgh, lay the foundations for the Prolog and LP of today.
- The “Year of Prolog” celebrates the 50th anniversary of these events.
Organizers: Association for Logic Programming and Prolog Heritage Association.
- Objectives:
 - ▶ Highlight the continuing significance of Prolog and LP for both symbolic, explainable AI, and computing more generally.
 - ▶ Inspire a new generation of students, by drawing their attention to the logic-based approach to computing.

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and others... do join in!

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 - ▶ What, 50 years?!? Half a century?!?!
 - ▶ Is Prolog therefore now 'old'? Is Prolog now irrelevant?

- Actually... continued interest:
 - ▶ Many *active implementations*, and *more appearing* continuously.
 - ▶ TIOBE index of programming languages shows Prolog:
 - In upper 10% of all languages tracked (270).
 - Stable, even somewhat upward trend since 2012.
 - One of only 13 languages that are tracked 'long term'.
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Early steps, major milestones

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See Kowalski (1988, 2013), Cohen (1988), VanRoy (1994), Colmerauer (1996), Gupta et al. (2001), vanEmden (2006), McJones's archive, etc.
- Anyway, some highlights:
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 - ▶ Green (1969): extend resolution to answer questions in FO-logic (QA3).
 - ▶ Colmerauer (1970): Q-systems.
 - ▶ Kowalski and Kuehner (1971): SL-resolution (focused search).
 - ▶ Boyer and Moore (1972): structure sharing.
 - Marseilles - Edinburgh collaboration (Colmerauer/Kowalski and teams).
 - Prolog! (1972–1973)
 - ▶ The competing “procedural” view of AI (e.g., Hewitt).
 - Prompted Kowalski to marry the procedural and logical views.
 - ▶ Edinburgh: DHD Warren, +Pereira(s)/Bowen/Byrd; later Lisbon.
 - Dec-10 Prolog



Early Prologs and main milestones (\approx up to ISO)

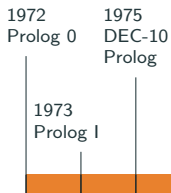
1972
Prolog 0

1973
Prolog I



- First *Prolog*(s): fundamental characteristics already there!

Early Prologs and main milestones (\approx up to ISO)



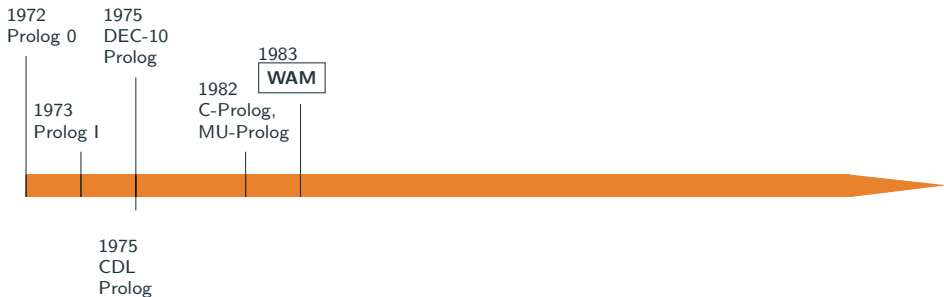
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- **The WAM:** portability + speed... and implementation beauty.

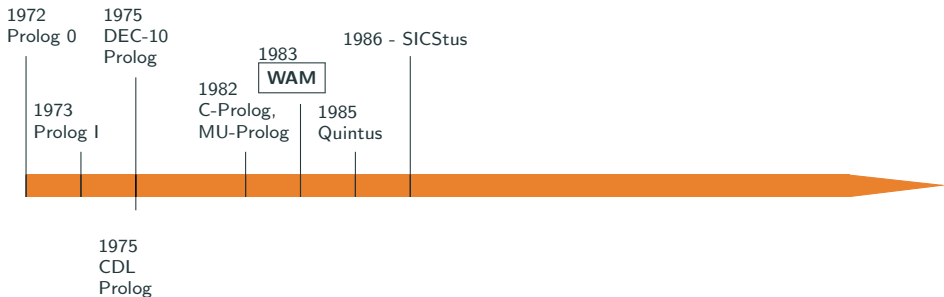
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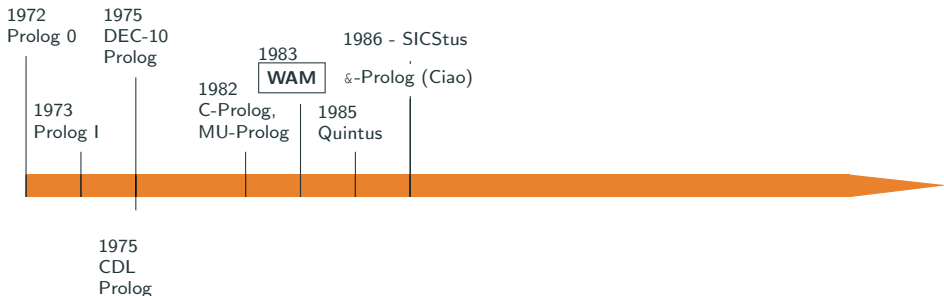
(FGCS → MCC → ECRC → ESPRIT → EU research programs, and others.)

Early Prologs and main milestones (\approx up to ISO)



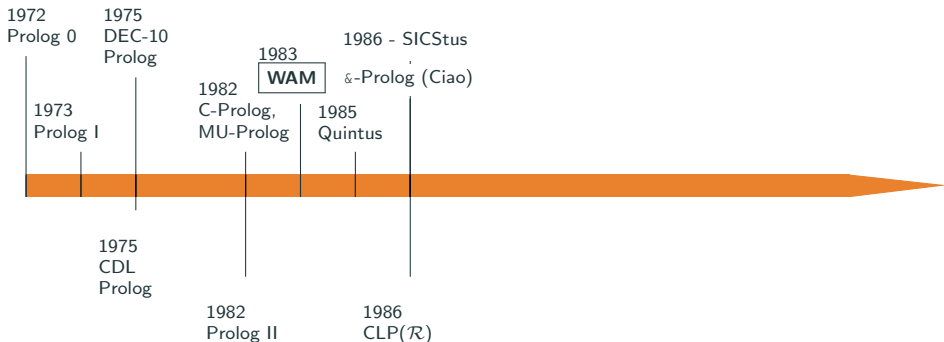
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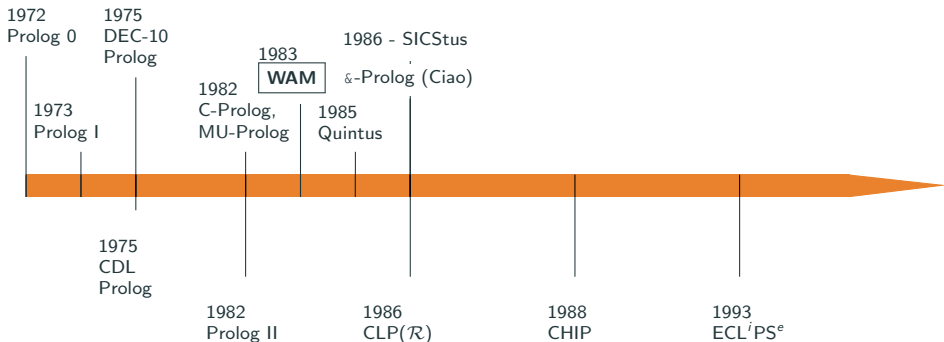
- WAM optimizations (Quintus, SICStus, BIM, YAP, ...), GC, ...
→ commercial/PD, dissemination, more performance.
- Or- and and-parallelism.
- *Global analysis* (abstract interpretation), P.Eval.; Aquarius, &-Prolog/Ciao.
(Independence/aliasing, modes, types, determinacy, sharing, non-failure, cost, ...)
First practical compiler(s) using abstract interpretation?
→ Performance (\approx imperative), auto-parallelization, real parallel speedups.

Early Prologs and main milestones (\approx up to ISO)



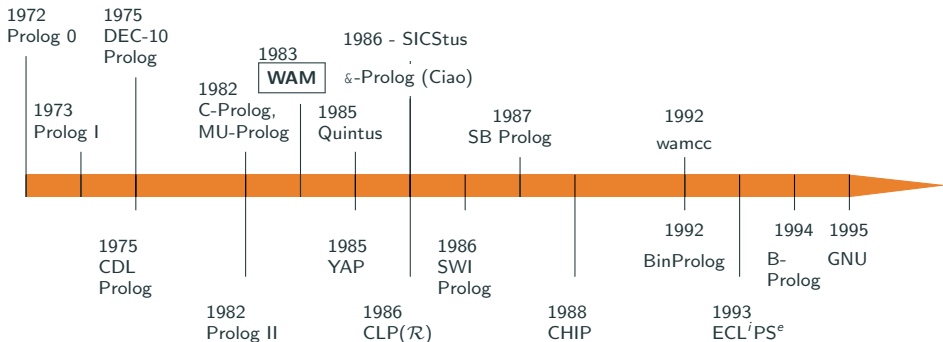
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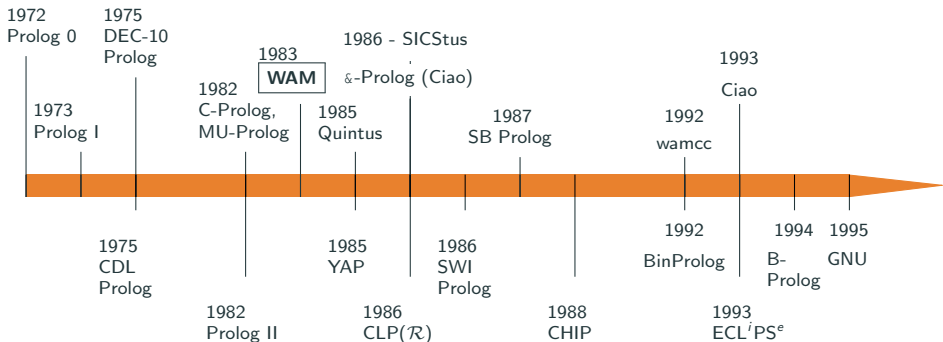
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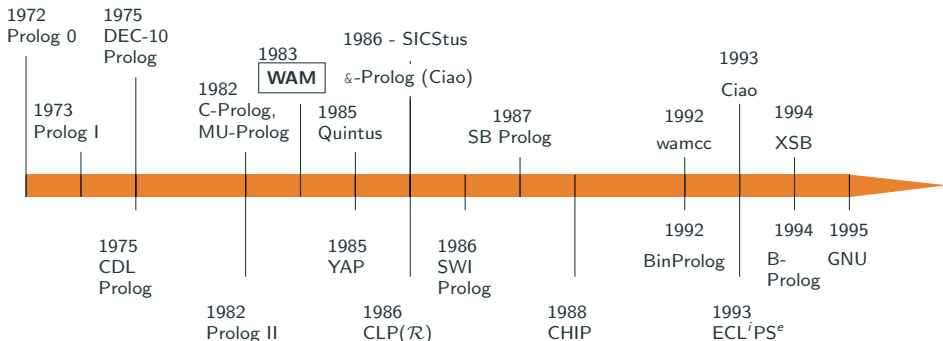
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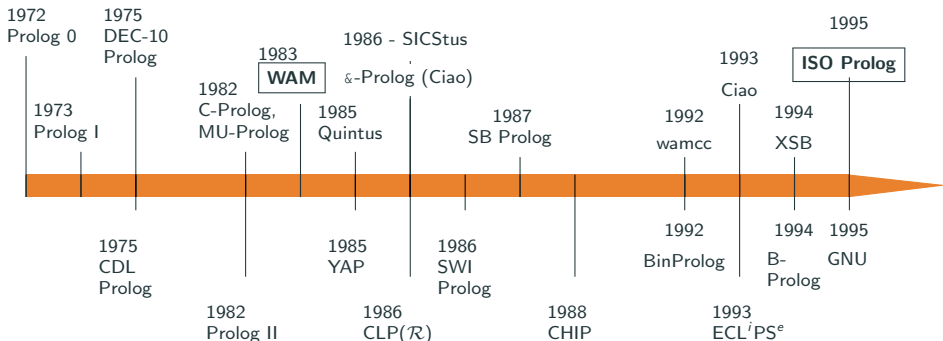
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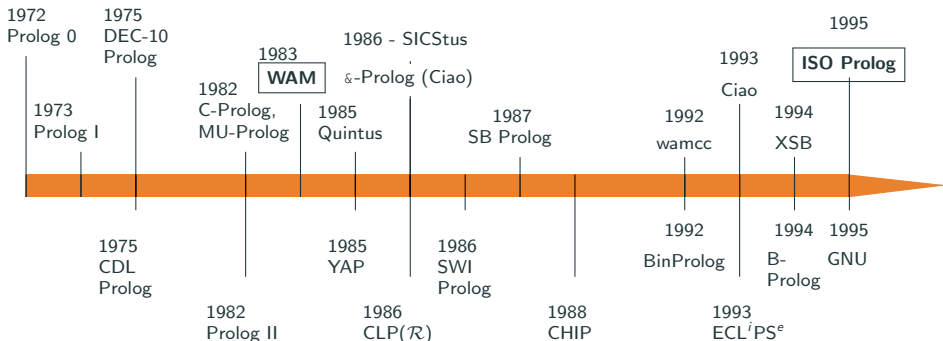
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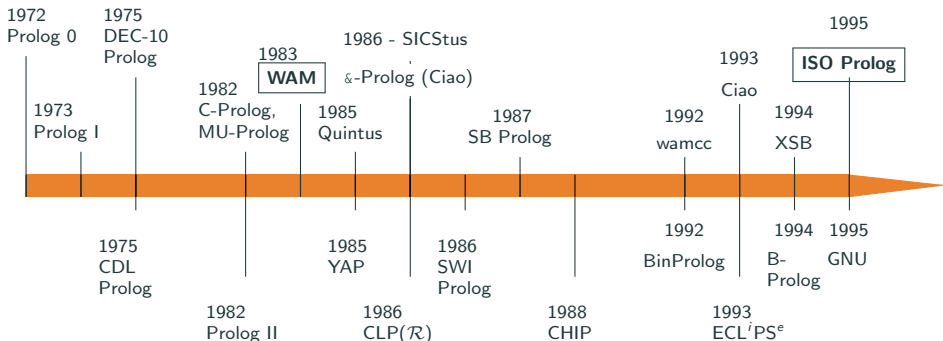
Early Prologs and main milestones (\approx up to ISO)



All this progressed in parallel with further advances in the theoretical underpinnings:

- Kowalski/van Emden (1976): linear res. for Horn clauses, no factoring rule, ...
- Clark (1978): correctness of NaF w.r.t. program completion.
- Reiter (1978): formalization of “Closed world assumption.”
- Minker, Gallaire, Cohen, Lassez/Jaffar/Maher, DHD Warren, Tamaki/Sato, DS Warren, ...

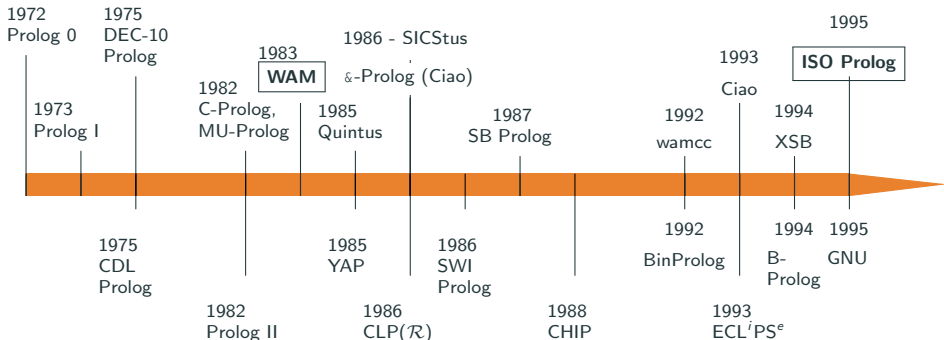
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After ISO – much additional evolution:

- Constraints in standard Prologs: “Opening the box” (attvars/CHR).
 - Learning (ILP), probabilistic.
 - *ASP* \rightsquigarrow Prolog-ASP combinations \rightsquigarrow *s(CASP)*.
 - Web embedding, playgrounds, notebooks.
- + applications of techniques to other languages,
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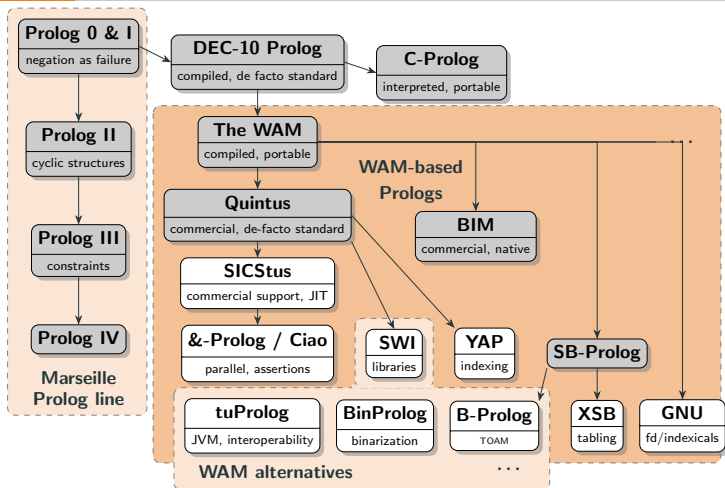
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Let's jump forward and take a look at the current state of things!

An overview of current systems

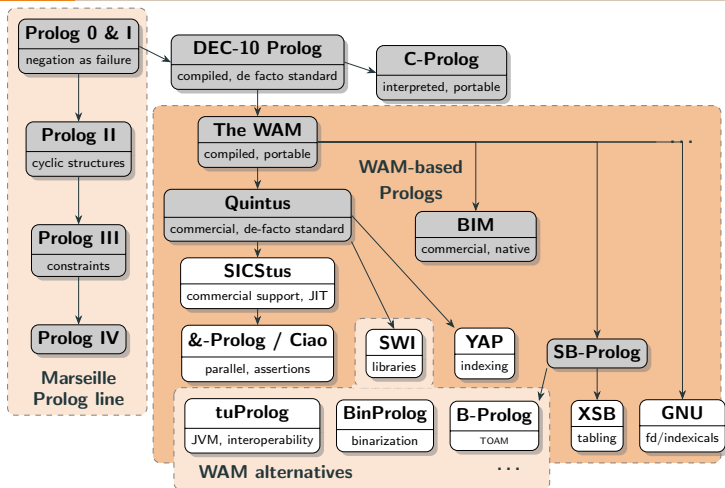
Prolog system heritage



White background: currently active/supported systems.
Lower legends: just some highlight(s) (see later).
Arrows: influences and inspiration.

Again, more missing!: ECLⁱPS^e, IBM, LIFE, Andorra-I, Scryer, Tau, ...

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Support status for selected features - I

System	Open Src.	Modules	Non-Std. Data Types	Foreign Language Interfaces
B-Prolog			arrays, sets, hashtables	C, Java
Ciao	✓	✓		C, Java, Python, JScript
ECLiPSe	✓	✓	arrays, strings	C, Java, Python, PHP
GNU Prolog	✓		arrays	C, Java, PHP
JIProlog	✓	✓		Java
SICStus		✓		C, Java, .NET, Tcl/Tk
SWI	✓	✓	dicts, strings	C, C++, Java
τ Prolog	✓	✓		JavaScript
tuProlog	✓		arrays	Java, .NET, Android, iOS
XSB	✓	✓		C, Java, PERL, Python
YAP	✓	✓		C, Python, R

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Support status for selected features - II

System	CLP	CHR	Tabling	Parallelism	Indexing	Coroutines
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Ciao	<i>FD, Q, R</i>	✓	✓	✓	FA, MA	✓
ECLiPSe	<i>FD, Q, R, Set</i>	✓		✓	most suitable	✓
GNU Prolog	<i>FD, B</i>				FA	
JIProlog					undocumented	
SICStus	<i>FD, B, Q, R</i>	✓			FA	✓
SWI	<i>FD, B, Q, R</i>	✓	✓	✓	MA, deep, JIT	✓
τ Prolog					undocumented	
tuProlog				✓	FA	
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System	Debugger	Global Vars.	Mutables	Testing	Types/Modes	s(CASP)
B-Prolog	trace	✓				
Ciao	trace / source	✓	✓	✓	✓	✓
ECLiPSe	trace	✓		✓		
GNU Prolog	trace	✓	✓			
JIProlog	trace					
SICStus	trace / source		✓	✓		
SWI	trace / graphical	✓	✓	✓		✓
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XSB	trace					
YAP	trace	✓				

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YAP	trace	✓				

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SICStus	trace / source		✓	✓		
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XSB	trace					
YAP	trace	✓				

Many other features!

- Auto-documentation, attributed variables, objects, enhanced expansions, playgrounds, ...

Summary (so far)

- Prolog systems have come a long way!
- ISO standard generally supported (with minor differences).
- *Basic* module system pretty compatible.
- A good number of commonly available features:
 - ▶ Constraints.
 - ▶ Multi-threading.
 - ▶ Tabling.
 - ▶ Coroutining.
 - ▶ ...

However,

- ▶ Interfaces and details often differ.
 - Can mostly be bridged (c.f., Paolo Moura's work), but a real nuisance.
- ▶ Some features (e.g., Types/modes/verification, s(CASP), ...) still in few systems.

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Influences on others

Influence in other languages within LP and its extensions

- Goedel, Mercury, Turbo-Prolog (static typing)
- λ -Prolog, Curry, Babel
- CP, GHC, Parlog, Erlang (committed choice)
- Datalog, ASP
- s(ASP) and s(CASP) (can also be seen as extensions)
- HyProlog, Co-inductive LP, ...
- Probabilistic LP
- LogTalk
- Picat
- CHR, CHR_G
- ...

Influence beyond LP

- Theorem proving technology.
- Java (abstract machine, specification, ...).
- Erlang.
- Many embeddings in other languages.
- Many others: C++, many compilers, ...
- Analyzers and verifiers for other languages.
- ...

Further analysis of current status and outlook

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Prolog weaknesses → and how to address them

- Learning curve, beginners can easily write programs that loop or consume a huge amount of resources → teach it well, use the right tools! (see later)
- Lack of static typing → but notable exceptions!
- Lack of data hiding → but notable exceptions!
- Lack of object orientation. → but notable exceptions!
- Packages: availability and management → improve compatibility.
- Limited support for embedded or app development → but notable exceptions!

- Syntactically different from “traditional” programming languages, not a mainstream language → offer alternative syntax?
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 - ▶ Big Data.
- New features and developments:
 - ▶ Probabilistic reasoning.
 - ▶ Embedding ASP and SAT or SMT solving, s(CASP) applications.
 - ▶ Opportunity still for performance gains (and we have the technology):
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 - ▶ Opportunity still for performance gains (and we have the technology):
 - Full-fledged JIT compiler.
 - Global optimization, partial evaluation ('provably correct refactoring').
 - Parallelism.
 - ▶ ...

Opportunities for Prolog

- New application areas, addressing societal challenges:
 - ▶ Neuro-Symbolic AI.
 - ▶ Explainable AI, verifiable AI.
 - ▶ Big Data.
- New features and developments:
 - ▶ Probabilistic reasoning.
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- Improve portability of existing features (cf., Prolog systems tables):
 - ▶ ISO, vs. Prolog Commons, vs. future initiatives,
 - ▶ Library infrastructure and conditional code,
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Types, modes, and other properties

(Some perspectives from the Ciao Prolog system)

Dynamic vs. Static languages – the classic dilemma

Dynamic languages (Prolog, Lisp/Scheme, Python, Javascript, ...)

- Dynamic checking of basic types, modes, and some other properties:
 - ▶ `..., A is B+C, ...`
B and C checked by `is/2` to be *instantiated* to `numexpr` at run time.
 - ▶ `..., arg(N,T,A), ...`
N checked by `arg/3` to be `nat` & $\leq \text{arity}(T)$ ("array bounds").
- Flexibility, compactness, rapid prototyping, scripting, ..., but
- ▶ Most errors only detected at run time.
 - ▶ Need to use tags (*boxing* of data) to identify type and mode, store arity, etc.

Static languages (ML, Haskell, Mercury, Gödel, ...)

- Compiler statically checks *types*.
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Solving the Dynamic vs. Static Dilemma

The Ciao Approach (mid 90s's!):

1. *Assertions* can be used to express types, modes, and many other properties.
 - ▶ But voluntary: provided up front, gradually, or not at all.
2. Then, *advanced program analysis* (abstract interpretation) is used to:
 - ▶ Verify the assertions:
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 - ▶ Achieve high performance:
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 - Unboxing, specialization, slicing, automatic parallelization, ...
3. Also, easily generate tests from assertions (this is (C)LP!).

- Provides the flexibility of dynamic languages, but with
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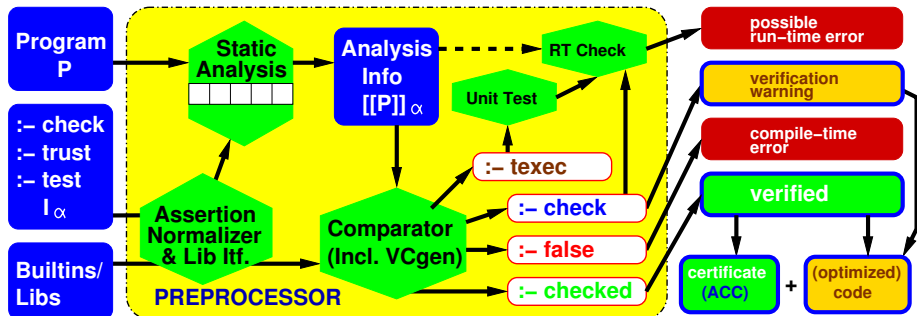
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The Ciao Integrated Approach to Specification, Debugging, Verification, Testing, and Optimization



Discussion: Comparison with Classical Types

"Traditional" Types	Ciao Assertion-based Model
"Properties" limited by decidability	Much more general property language
May need to limit prog. lang.	No need to limit prog. lang.
"Untypable" programs rejected	Run-time checks introduced
(Almost) Decidable	Decidable + Undecidable (approximated)
Expressed in a different language	Expressed in the source language
Types must be defined	Types can be defined or inferred
Assertions are only of type "check"	"check", "trust", ...
Type signatures & assertions different	Type signatures <i>are</i> assertions

- But quite popular now: gradual typing, Racket, liquid Haskell, etc.
- Some key issues:
 - Safe / Sound approximation*
 - Abstract Interpretation*
 - Suitable assertion language*
 - Powerful abstract domains*
- Works best when properties and assertions can be expressed in the source language (i.e., source lang. supports *predicates, constraints*).

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Demo! (See slides at the end.)

Teaching (and preaching) Prolog

On teaching (and preaching) Prolog

- Prolog / LP / CLP *must* be taught in CS programs,
 - ▶ A CS graduate is simply not complete without knowledge of Prolog.
(and maybe also in other majors and maybe in schools –cf. Prolog Year?)
- But it has to be done right!
 - ▶ The standard 'programming paradigms' approach can be counter-productive.
 - ▶ Simply cannot be done in a couple of weeks emulating Prolog in Scheme.
 - What to do if that is the only slot available?
- On the way *dispel unfounded myths* about the language, and show how many of the shortcomings of early Prologs have been *addressed over the years*.

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On teaching (and preaching) Prolog

- “Prolog gets into infinite loops.”

This is true –in fact, of any programming language or proof system. However, it is likely to discourage beginners if not explained well:

- ▶ Use a system that can *alternatively and selectively* run in breadth-first, iterative deepening, tabling, etc.
- ▶ Start by running all predicates, e.g., breadth-first – everything works!
- ▶ Then, explain the shape of the tree (solutions at finite depth, possible infinite failures, etc.), and thus why breadth-first works, and why depth-first sometimes may not.
- ▶ Do relate it to the *halting problem*: no-one (Prolog, logic, nor other Turing-complete prog. language) can solve that (but tabling helps).
- ▶ Discuss advantages and disadvantages of search rules (time, memory). Motivate the choices made for Prolog benchmarking actual executions.

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- ▶ Do relate it to the *halting problem*: no-one (Prolog, logic, nor other Turing-complete prog. language) can solve that (but tabling helps).
- ▶ Discuss advantages and disadvantages of search rules (time, memory). Motivate the choices made for Prolog benchmarking actual executions.

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- “Prolog gets into infinite loops.”

This is true –in fact, of any programming language or proof system.

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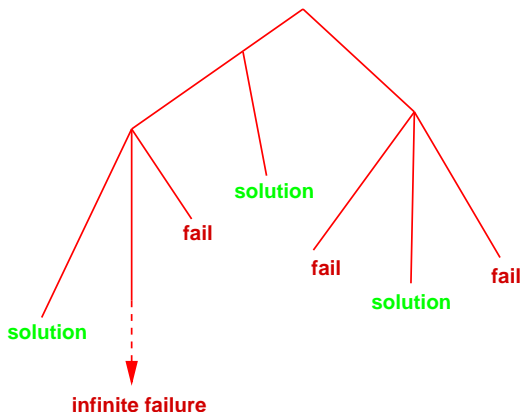
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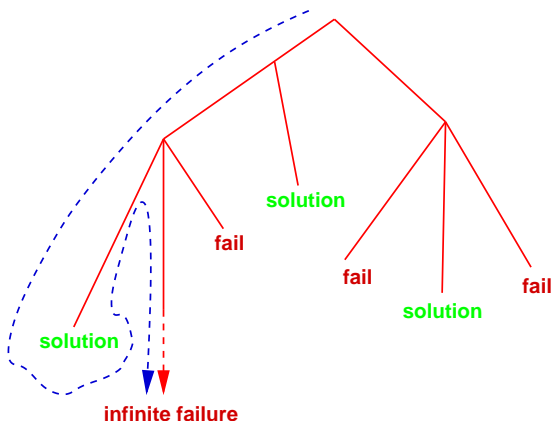
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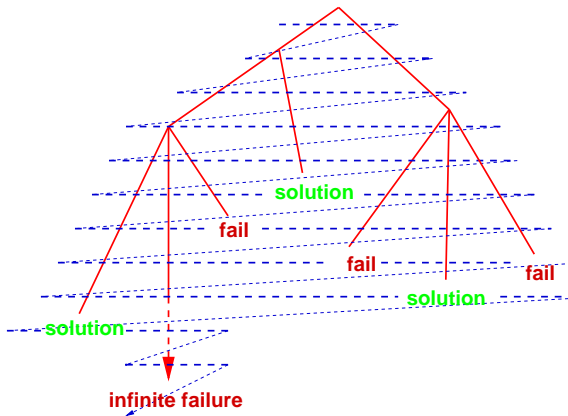
Characterization of the search tree



Depth-First Search



Breadth-First Search



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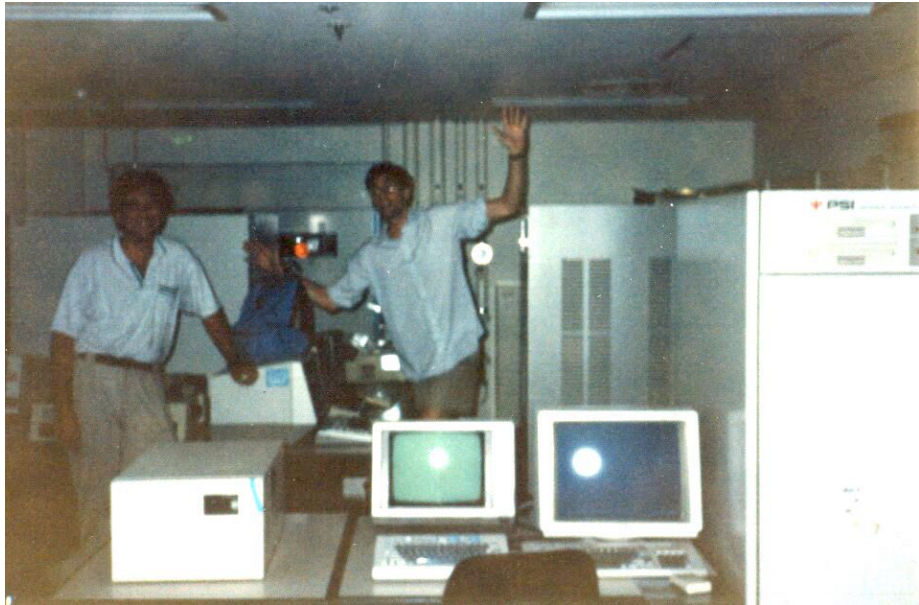
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Personal Sequential Inference –PSI– machine (Prolog machine) in FGCS ICOT's basement (the large refrigerator-size box on the right).

On teaching (and preaching) Prolog

- Do show the beauty:
 - ▶ Explain “Green’s dream,” discuss for what logics we have effective deduction procedures, justify the choice of FO and semi-decidability, SLD-resolution → classical LP (Kowalski/Colmerauer).
 - ▶ Show how logic programs are both logical theories (with declarative meaning) and procedural programs that can be debugged, followed step by step, etc.
 - An operational (in addition to declarative) semantics is a requirement in the language (vs., e.g., Goedel) and we do need to teach it.
 - Otherwise not a programming language, just specification/KR – Prolog is both.
 - How otherwise to reason about complexity, memory consumption, etc.? To say that these things don’t matter does not make sense in PL.
 - ▶ Show with examples (and benchmarking them) how you can go from executable specifications to efficient algorithms gradually, and as needed.
 - ▶ Show how unification is also a device for *constructing and matching complex data structures with (declarative) pointers*. Show it in the top level, giving “the data structures class.”
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 - ▶ Some parts of it can result from the Year of Prolog efforts.

Demo slides for the part on:

Types, modes, and other properties

(Some perspectives from the Ciao Prolog system)

Example: qsort

Ciao warns that it cannot verify that the call to `=</2` will not generate a run-time error (assertion is in library!):

```
»:- module(., [qsort/2], [assertions, nativeprops, (nmodes)]).

qsort([], []).
qsort([First|Rest], Result) :-
    partition(Rest, First, Sm, Lg),
    qsort(Sm, SmS),
    qsort(Lg, LgS),
    append(SmS, [First|LgS], Result).

partition([], _, [], []).
»partition([X|Y].F, [X|Y1], Y2) :-
    X =< F,
    partition(Y, F, Y1, Y2).
»partition([X|Y].F, Y1, [X|Y2]) :-
    X > F,
    partition(Y, F, Y1, Y2).

append([], Xs, Xs).
append([X|Xs], Ys, [X|Zs]) :-
    append(Xs, Ys, Zs).
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partition([],_,[],[]).
»partition([X|Y],F,[X|Y1],Y2) :-
    > At literal 1 could not verify assertion:
    partition(Y,F,Y1,Y2).
»partition([X|Y],F,Y1,[X|Y2]) :-
    X > F,
    partition(Y,F,Y1,Y2).

append([],Xs,Xs).
append([X|Xs],Ys,[X|Zs]) :-
    append(Xs,Ys,Zs).
```

Example: qsort

Adding useful entry information Ciao can infer that $=</2$ is called correctly, and no warnings are flagged (this would normally be obtained from analysis of caller to this module):

```
:- module(_, [qsort/2], [assertions, nativeprops, .(nmodes)]).  
  
:- pred qsort(+list(num), _).  
  
qsort([], []).  
qsort([First|Rest], Result) :-  
    partition(Rest, First, Sm, Lg),  
    qsort(Sm, SmS),  
    qsort(Lg, LgS),  
    append(SmS, [First|LgS], Result).  
  
partition([], _, [], []).  
partition([X|Y], F, [X|Y1], Y2) :-  
    X =< F,  
    partition(Y, F, Y1, Y2).  
partition([X|Y], F, Y1, [X|Y2]) :-  
    X > F,  
    partition(Y, F, Y1, Y2).  
  
append([], Xs, Xs).  
append([X|Xs], Ys, [X|Zs]) :-  
    append(Xs, Ys, Zs).
```


Example: qsort

We add some more assertions... :

```
:- pred qsort(+list(num),-list(num)) + is_det.

qsort([], []).
qsort([First|Rest],Result) :-
    partition(Rest,First,Sm,Lg),
    qsort(Sm,SmS),
    qsort(Lg,LgS),
    append(SmS,[First|LgS],Result).

:- pred partition(+list(num),+num,-list(num),-list(num)) + (is_det,not_fails).

partition([],_,[],[]).
partition([X|Y],F,[X|Y1],Y2) :-
    X <= F,
    partition(Y,F,Y1,Y2).
partition([X|Y],F,Y1,[X|Y2]) :-
    X > F,
    partition(Y,F,Y1,Y2).

:- pred append(+list(num),+list(num),-list(num)) + is_det.

append([],Xs,Xs).
append([X|Xs],Ys,[X|Zs]) :-
    append(Xs,Ys,Zs).
```

Example: qsort

...and they get verified by Ciao:

```
»:- pred qsort(+list(num),-list(num)) + is_det.

qsort([], []).
qsort([First|Rest],Result) :-
    partition(Rest,First,Sm,Lg),
    qsort(Sm,SmS),
    qsort(Lg,LgS),
    append(SmS,[First|LgS],Result).

»:- pred partition(+list(num),+num,-list(num),-list(num)) + (is_det,not fails).

partition([],_,[],[]).
partition([X|Y],F,[X|Y1],Y2) :-
    X <= F,
    partition(Y,F,Y1,Y2).
partition([X|Y],F,Y1,[X|Y2]) :-
    X > F,
    partition(Y,F,Y1,Y2).

»:- pred append(+list(num),+list(num),-list(num)) + is_det.

append([],Xs,Xs).
append([X|Xs],Ys,[X|Zs]) :-
    append(Xs,Ys,Zs).
```

Example: qsort

...and they get verified by Ciao:

```
»:- pred qsort(+list(num),-list(num)) + is det.
```

```
qsort([], []).
```

```
qsort([First|Rest],Result) :-  
    partition(Rest,First,Sm,Lg),  
    qsort(Sm,SmS),  
    qsort(Lg,LgS),  
    append(SmS,[First|LgS],Result).
```

```
»- pred partition(+list(num),+num,-list(num),-list(num)) + (is det,not fails).
```

```
> Verified assertion:
```

```
:- check comp partition(A,B,C,D)  
   : ( list(num,A), num(B) )  
   + ( is_det, not_fails ).
```

```
> Verified assertion:
```

```
:- check success partition(A,B,C,D)  
   : ( list(num,A), num(B) )  
   => ( list(num,C), list(num,D) ).
```

```
»:- pred append(+list(num),+list(num),-list(num)) + is det.
```

```
append([],Xs,Xs).
```

```
append([X|Xs],Ys,[X|Zs]) :-  
    append(Xs,Ys,Zs).
```

Example: qsort

If we replace `=</2` with `</2` Ciao warns that `partition/3` can fail (cannot prove `not_fails`):

```
»:- pred qsort(+list(num),-list(num)) + is_det.

qsort([], []).
qsort([First|Rest],Result) :-
    partition(Rest,First,Sm,Lg),
    qsort(Sm,SmS),
    qsort(Lg,LgS),
    append(SmS,[First|LgS],Result).

»:- pred partition(+list(num),+num,-list(num),-list(num)) + (is_det,not_fails).

partition([],_,[],[]).
partition([X|Y],F,[X|Y1],Y2) :-
    X < F,
    partition(Y,F,Y1,Y2).
partition([X|Y],F,Y1,[X|Y2]) :-
    X > F,
    partition(Y,F,Y1,Y2).

»:- pred append(+list(num),+list(num),-list(num)) + is_det.

append([],Xs,Xs).
append([X|Xs],Ys,[X|Zs]) :-
    append(Xs,Ys,Zs).
```

Example: qsort

If we replace $\geq/2$ with $>/2$ Ciao warns that `partition/3` is not deterministic (cannot prove `is_det`):

```
»:- pred qsort(+list(num),-list(num)) + is_det.

qsort([], []).
qsort([First|Rest],Result) :-
    partition(Rest,First,Sm,Lg),
    qsort(Sm,SmS),
    qsort(Lg,LgS),
    append(SmS,[First|LgS],Result).

»:- pred partition(+list(num),+num,-list(num),-list(num)) + (is_det,not fails).

partition([],_,[],[]).
partition([X|Y],F,[X|Y1],Y2) :-
    X <= F,
    partition(Y,F,Y1,Y2).
partition([X|Y],F,Y1,[X|Y2]) :-
    X >= F,
    partition(Y,F,Y1,Y2).

»:- pred append(+list(num),+list(num),-list(num)) + is_det.

append([],Xs,Xs).
append([X|Xs],Ys,[X|Zs]) :-
    append(Xs,Ys,Zs).
```

Example: nrev (using the functional syntax package)

An example with more complex properties, a cost error is flagged:

```
»:- module( , [nrev/2], [assertions,fsyntax,nativeprops]).  
»:- pred nrev(A,B) : {list, ground} * var => list(B)  
  + ( not_fails, is_det, steps_o( length(A) ) ).  
  
nrev( [] )      := [].  
nrev( [H|L] )  := ~conc( ~nrev(L), [H] ).  
  
:- pred conc(A,B,C) + ( terminates, is_det, steps_o(length(A)) ).  
  
conc( [], L ) := L.  
conc( [H|L], K ) := [ H | ~conc(L,K) ].
```

Example: nrev

Ciao reminds us that `nrev/2` is of course quadratic, not linear:

```
»:- module( ., [nrev/2], [assertions,fsyntax,nativeprops]).
»:- pred nrev(A,B) : {list, ground} * var => list(B)
  > False assertion:
  :- check comp nrev(A,B)
     : ( list(A), ground(A), var(B) )
     + ( not_fails, is_det, steps_o(length(A)) ).
  because the comp field is incompatible with inferred comp:
  [generic_comp] covered,is_det,mut_exclusive,not_fails,steps_lb(0.5*exp(length(A)
  » ,2)+1.5*length(A)+1),steps_ub(0.5*exp(length(A),2)+1.5*length(A)+1)
  > Verified assertion:
  :- check calls nrev(A,B)
     : ( list(A), ground(A), var(B) ).
  > Verified assertion:
  :- check success nrev(A,B)
     : ( list(A), ground(A), var(B) )
     => list(B).
```

Example: nrev

With the cost expression fixed all properties are now verified:

```
>>:- module( , [nrev/2], [assertions,fsyntax,nativeprops]).  
  
>>:- pred nrev(A,B) : {list,ground} * var => list(B)  
    + ( not_fails, is_det, steps_o( exp(length(A),2) ) ).  
  
nrev( [] )      := [].  
nrev( [H|L] )  := ~conc( ~nrev(L),[H] ).  
  
>>:- pred conc(A,B,C) + ( terminates, is_det, steps_o(length(A)) ).  
  
conc( [], L ) := L.  
conc( [H|L], K ) := [ H | ~conc(L,K) ].
```


Example: nrev

If we change the assertion for `conc/3` from complexity order (`_o`) to upper bound (`_ub`) then Ciao flags that `length(A)` is not a correct upper bound:

```
»:- module( ., [nrev/2], [assertions,fsyntax,nativeprops]).  
»:- pred nrev(A,B) : {list,ground} * var => list(B)  
  + ( not_fails, is_det, steps_o( exp(length(A),2) ) ).  
  
nrev( [] )      := [].  
nrev( [H|L] )  := ~conc( ~nrev(L),[H] ).  
  
»:- pred conc(A,B,C) + ( terminates, is_det, steps_ub(length(A)) ).  
  
conc( [], L ) := L.  
conc( [H|L], K ) := [ H | ~conc(L,K) ].
```

Example: nrev

If we change the assertion for `conc/3` from complexity order (`_o`) to upper bound (`_ub`) then Ciao flags that `length(A)` is not a correct upper bound:

```
⌘:- module( ., [nrev/2], [assertions,fsyntax,nativeprops]).
⌘:- pred nrev(A,B) : {list,ground} * var => list(B)
    + ( not_fails, is_det, steps_o( exp(length(A),2) ) ).

nrev( [] )      := [].
nrev( [H|L] )  := ~conc( ~nrev(L),[H] ).

⌘- pred conc(A,B,C) + ( terminates, is_det, steps_ub(length(A)) ).
> False assertion:
:- check comp conc(A,B,C)
    + ( terminates, is_det, steps_ub(length(A)) ).
because the comp field is incompatible with inferred comp:
[generic_comp] covered,is_det,mult_exclusive,not_fails,steps_lb(length(A)+1),step
s_ub(length(A)+1)
> Verified assertion:
:- check calls conc(A,B,C).
```

Example: nrev

With the cost expression fixed all properties are now verified:

```
>:- module( , [nrev/2], [assertions, fsyntax, nativeprops]).  
>:- pred nrev(A,B) : {list, ground} * var => list(B)  
    + ( not_fails, is_det, steps_o( exp(length(A),2) ) ).  
  
nrev( [] )      := [].  
nrev( [H|L] )  := ~conc( ~nrev(L), [H] ).  
  
>:- pred conc(A,B,C) + ( terminates, is_det, steps_ub(length(A)+1) ).  
  
conc( [], L ) := L.  
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  > Verified assertion:  
  :- check calls conc(A,B,C).  
  > Verified assertion:  
  :- check comp conc(A,B,C)  
    + ( terminates, is_det, steps_ub(length(A)+1) ).
```