

WHY VALIDATE STATIC ANALYZERS?

Static analysis tools are crucial in modern software development: verification, optimization, etc.

But building analyzers is **hard**:

- ▶ **Complex, large systems**
- ▶ **Prone to subtle bugs**
- ▶ **Used in critical tasks** needing trustworthy results

Validating analyzers is also difficult!

- ▶ **Formal methods** are hard to apply directly
- ▶ **Specifications** are often missing or not complete
- ▶ The *who checks the checker?* problem rises

In practice, extensive **testing** is the **most realistic** option.

- ▶ But, unit-tests miss integration bugs
- ▶ It's hard to define **testing oracles**
- ▶ Generation of **complex data structures** with hard-to-test conditions is challenging

HOW? *Checkification* Algorithm

1 Source program or set of benchmarks

tree.pl

```
:- module(tree,_,[assertions]).

:- pred insert(+nnegint,+tree,-).

insert(X,empty,tree(empty,X,empty)).
insert(X,tree(LC,X,RC),tree(LC,X,RC)).
insert(X,tree(LC,Y,RC),tree(LC_p,Y,RC):-
    X<Y,
    insert(X,tree(X,LC,LC_p)).
insert(X,tree(LC,Y,RC),tree(LC,Y,RC_p):-
    X>Y,
    insert(X,tree(X,RC,RC_p)).

:- pred belongs(+nnegint,+tree).
...
:- pred root(+tree,-).
...
```

2 Program is analyzed with domain \mathcal{D}

tree_analysis.pl

```
:- module(tree_analysis,_,[assertions]).

:- true insert(N,T1,T2)
    : (nnegint(N),tree(T1),var(T2))
    => (nnegint(N),tree(T1),non_empty_tree(T2)).

insert(X,empty,tree(empty,X,empty)).
insert(X,tree(LC,X,RC),tree(LC,X,RC)).
insert(X,tree(LC,Y,RC),tree(LC_p,Y,RC):-
    true(nnegint(X),tree(LC),nnegint(Y),tree(RC)),
    X<Y,
    true(...),
    insert(X,tree(X,LC,LC_p)),
    true(...).
insert(X,tree(LC,Y,RC),tree(LC,Y,RC_p):-
    true(...),
    X>Y,
    true(...),
    insert(X,tree(X,RC,RC_p)),
    true(...).
...
```

3 *true* status is replaced with *check*

tree_check.pl

```
:- module(tree_check,_,[assertions,rtchecks]).

:- check insert(N,T1,T2)
    : (nnegint(N),tree(T1),var(T2))
    => (nnegint(N),tree(T1),non_empty_tree(T2)).

insert(X,empty,tree(empty,X,empty)).
insert(X,tree(LC,X,RC),tree(LC,X,RC)).
insert(X,tree(LC,Y,RC),tree(LC_p,Y,RC):-
    check(nnegint(X),tree(LC),nnegint(Y),tree(RC)),
    X<Y,
    check(...),
    insert(X,tree(X,LC,LC_p)),
    check(...).
insert(X,tree(LC,Y,RC),tree(LC,Y,RC_p):-
    check(...),
    X>Y,
    check(...),
    insert(X,tree(X,RC,RC_p)),
    check(...).
...
```

4 Generation of tests

Use **assertion preconditions** as **generators** of valid test inputs.

→ Prolog's declarative nature lets properties be expressed as predicates.

Traditional Prolog uses **depth-first** search: **efficient** but **incomplete**.

Classic Ciao Prolog **allows other search strategies**.

💡 **PROPOSED SOLUTION**

A **new** mechanism to execute predicates under non-standard search rules (**breadth-first**, **id**, **random**, **guided**, ...) as well to explore **diverse inputs**.

✓ Not rewriting predicates, but running them with **alternative execution strategies**.

✓ More expressive specifications.

```
:- search_rule(tree/1,bf).

:- prop tree/1+regtype.
tree(empty).
tree(tree(LC,N,RC)):-
    tree(LC),
    gen([sr(df)],(nnegint(N)))
    tree(RC).

% Check if tree T is sorted
sorted_tree(T):- ...

% Constrains the sum of all node values in
% tree T to equal N
tsum(T,N):- ...

:- check insert(N,T1,T2)
    : (nnegint(N),tree(T1),var(T2))
    => (nnegint(N),tree(T1),non_empty_tree(T2)).
    + gen([sorted_tree(T1),tsum(T1,10)]).
```

✓ Facilitates **automatic + user-guided** test case generation.

✓ Pushes further towards general-purpose **flexible search** in (C)LP.

E.g., test cases `insert(N,T1,T2)` where **N** is instantiated to a non-negative integer, **T1** is a tree, and **T2** is a free variable.

But, in the `insert/3` specification, the generator is further guided by auxiliary properties: `sorted_tree/1` and `tsum/2`, which restrict the structure and content of the input trees.

- Test 1: `insert(2,tree(empty,10,empty),T)`
- Test 2: `insert(5,tree(empty,2,tree(empty,8,empty)),T)`
- ...
- Test n: `insert(8,tree(empty,1,tree(empty,4,tree(empty,5,empty))),T)`

If any of these **test cases** produce a **run-time error** then there is a **bug**!

I.e., if a run-time check reports a violation, then the analyzer must have inferred the assertion incorrectly, revealing a bug in the analyzer.

Current experimental outcomes

Many **applications**, depending on which parts of the system are **trusted**:

- + Debugging **Abstract Domains**.
- + Testing the **Abstract Interpretation Engine**.
Testing less trusted fixpoints and options (e.g., incremental analysis).
- + Debugging **trust assertions** and custom transfer functions
- + Testing the **overall consistency** of the **framework**.
E.g., when semantics is underspecified, check at least that runtime and static semantics agree.
- + Integration Testing of the Analyzer.
- + Testing **external** or **third party solvers** (e.g., PPL).

Analyzed programs with increasing levels of complexity:

- Success in finding **known bugs** or unsupported features in old versions (e.g., rational terms, attributed variables).
- Actual analysis bugs found, mainly in less mature domains.
- Some **inconsistencies** found in the **framework** (e.g., in interpretation of native properties by analyzer and runtime-checks).
- Some bugs in other components found and fixed (e.g., the analysis output).
- **Reasonable overhead**, with test execution time < 60s.

Full paper?

